



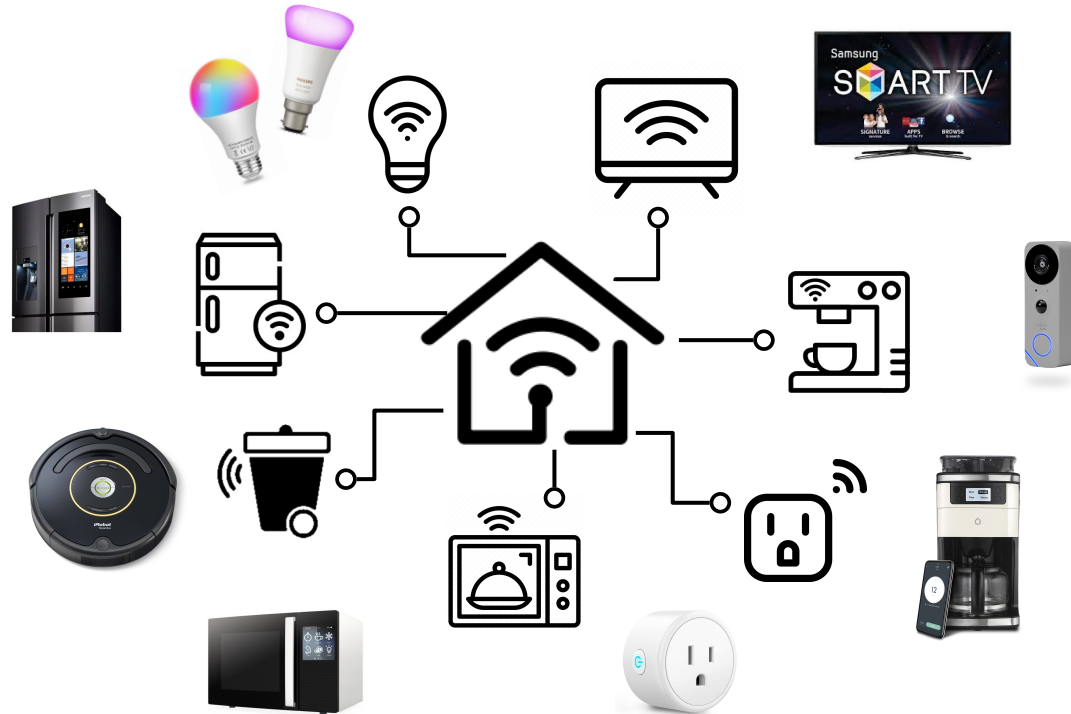
# Home, SafeHome: Smart Home Reliability with Visibility and Atomicity

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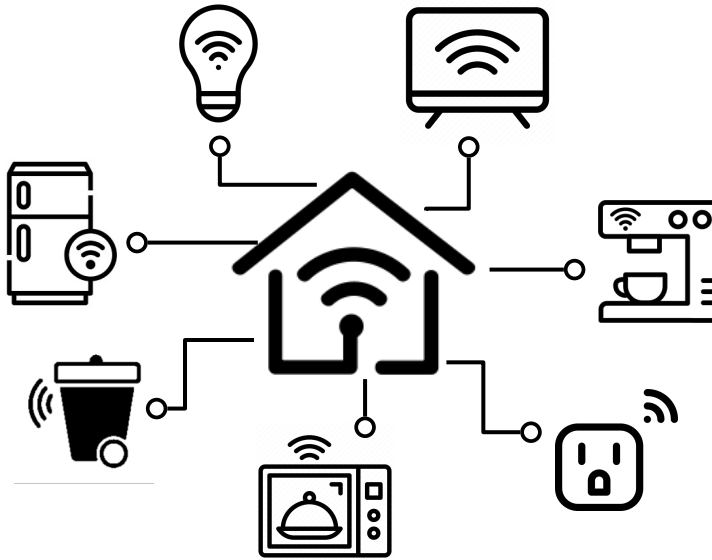


# Smart Home World



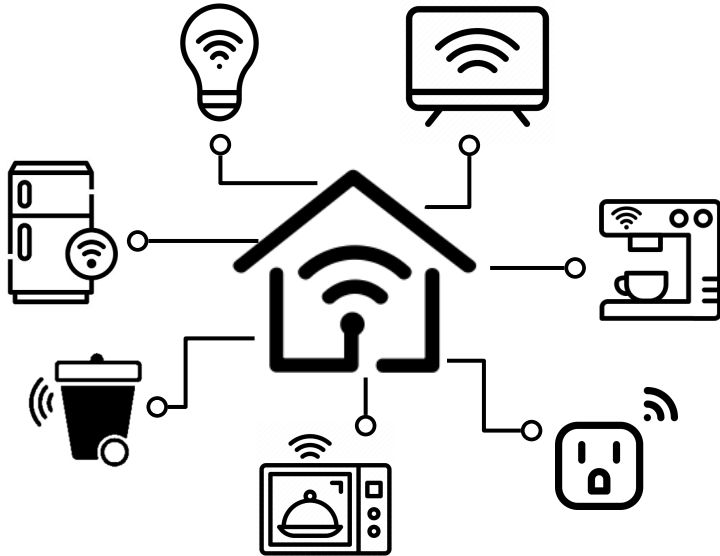
# Smart Home World

*Smart Device:* 1) connected to other devices via wireless protocols  
2) controlled by home automation systems



# Smart Home World

*Smart Device:* 1) connected to other devices via wireless protocols  
2) controlled by home automation systems



*“Humans need to control their lives, not control devices.”*

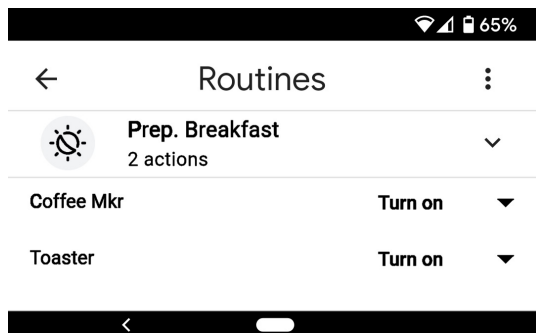
-- Davidoff et al, UbiComp'06

# Smart Home World (Cont.)

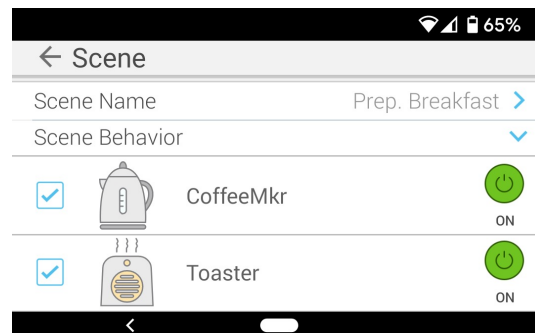
How people control smart home?

- by *Command*  
e.g. {Make an espresso}
- by *Routine*: a sequence of commands  
e.g. Prep. Breakfast = {Make an espresso; make a pancake}

Current systems execute Best-Effort!



Routine in Google Home



Routine in Kasa (TP-Link)



# Two Natural Expectations from Users

- Execute everything in a routine – *Atomicity*
  - All commands in the routine need to finish successfully, or none do
- When conflicts happen, people hope routines to execute one after another – *Isolation / Serial Equivalence*



*Poorly supported in  
current systems!*

*\*Routines are common to be long running, e.g. trash-out routine.*



# SafeHome

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- Home Automation System that can
  - Support *long running* routines
  - Properly *isolate* concurrent routines (providing *serial equivalence*)
  - Ensure routine execution *atomicity*
- Key challenge: Actions are visible to users
- Methodology:
  - Four *Visibility Models* (Spectrum for user choices)
  - *Lock-based* mechanism with *leasing* design



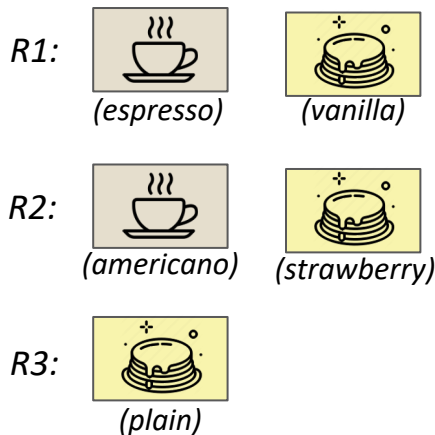
# Visibility Models

## Four Visibility Models:

- *Weak, Eventual, Partitioned Strict, Global Strict*

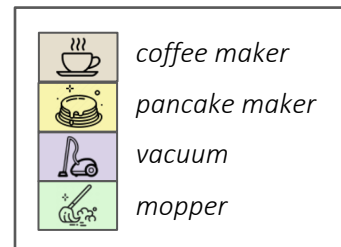
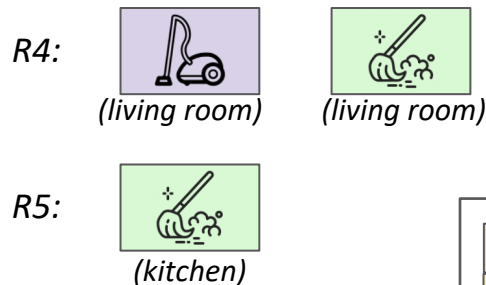
Example Scenarios: 5 routines are initiated *simultaneously* on 4 devices

3 Routines Initiated by User:  
Coffee Maker Pancake maker



&

2 Routines triggered by other sensors:  
Vacuum Mopper





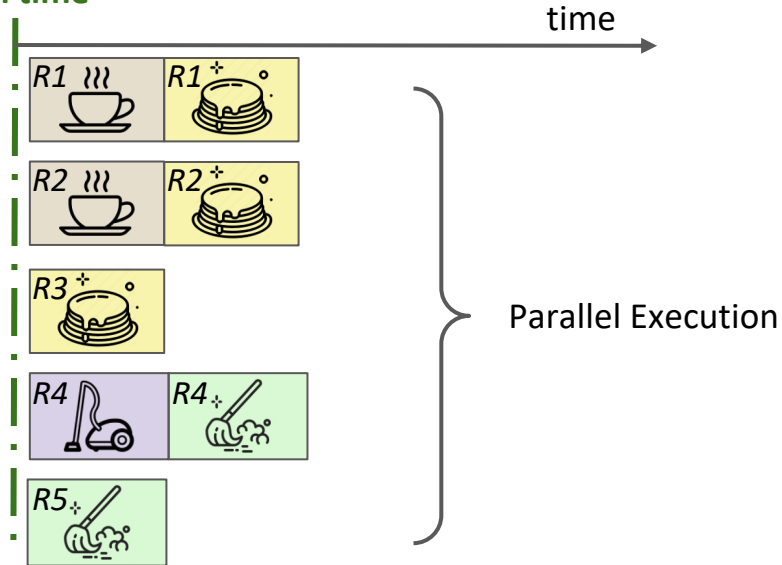
# Weak Visibility (WV) Model -- Status Quo

Strategy:





- Execute routine immediately when triggered

Finish in **2** time units

Insertion time



Two commands send simultaneously to one device may cause **errors**.

	coffee maker
	pancake maker
	vacuum
	mopper

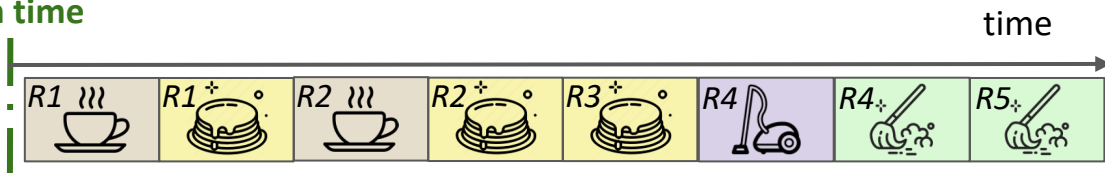
# Global Strict Visibility (GSV) Model

Strategy:





- Execute at most one routine at a time

Finish in 8 time units

Insertion time



- Strongest Visibility Model
- Example Usage: resource constrained environment:
  - e.g. 1000-watt max supply < coffee maker 600W + pancake maker: 600W

	coffee maker
	pancake maker
	vacuum
	mopper

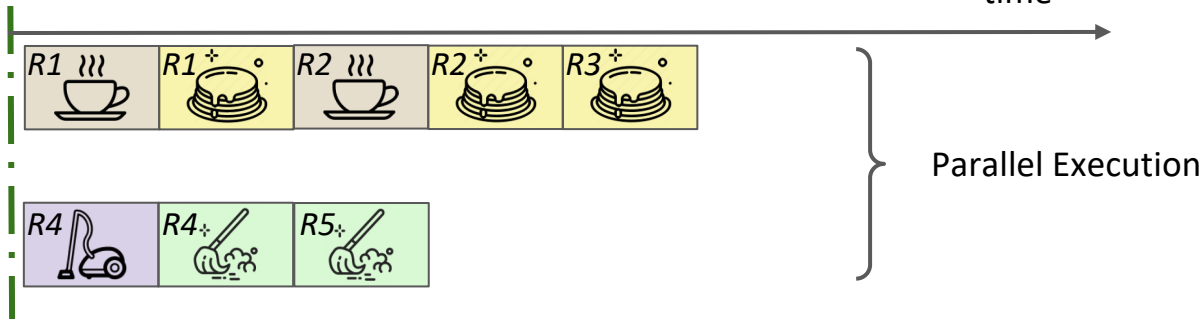
# Partitioned Strict Visibility (PSV) Model

Strategy:

- Routines touching disjoint devices do not block each other

Finish in **5** time units

Insertion time



- Useful when routines need to execute without interference through duration.
- Might still takes long with long running routines.



# Eventual Visibility (EV) Model

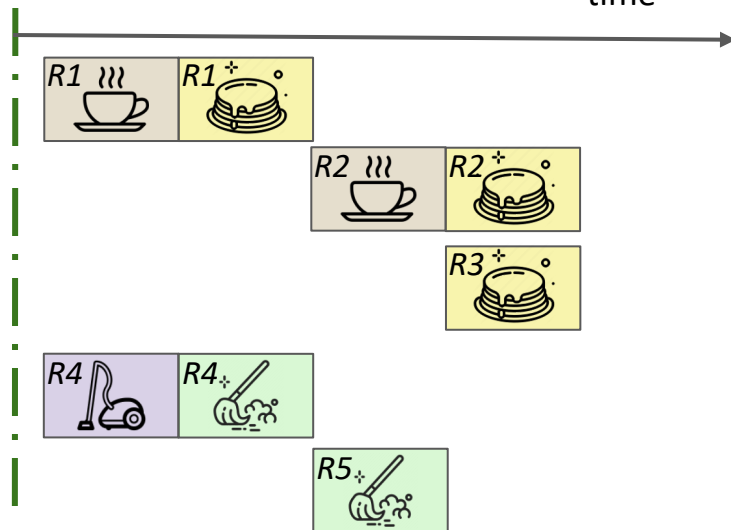
Strategy:

- Routines can concurrently execute *without violating some serial order*.

Finish in 3 time units

Insertion time

time



Parallel Execution

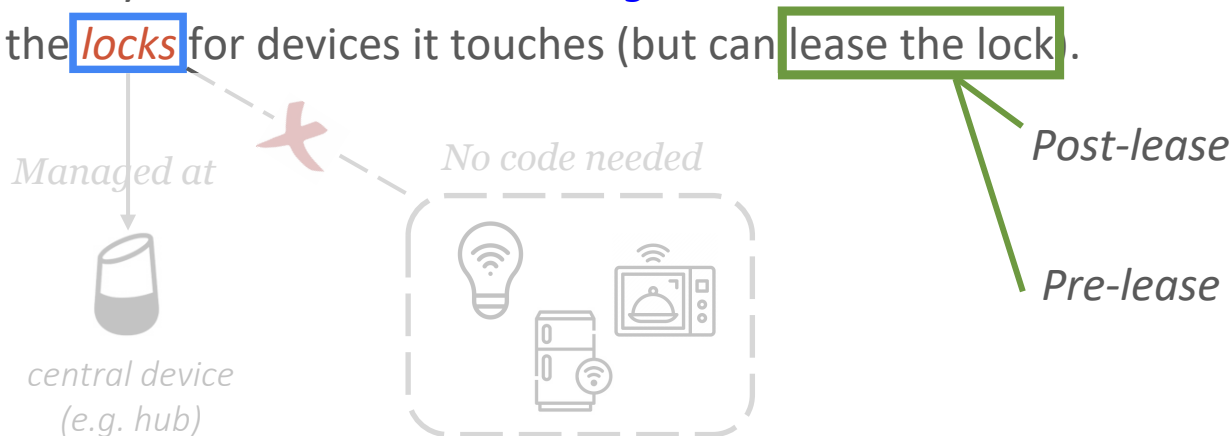
Equivalent end state to:  
 $R3 \rightarrow R1 \rightarrow R2 \rightarrow R5 \rightarrow R4$



# Eventual Visibility (EV) Model

## Strategy:

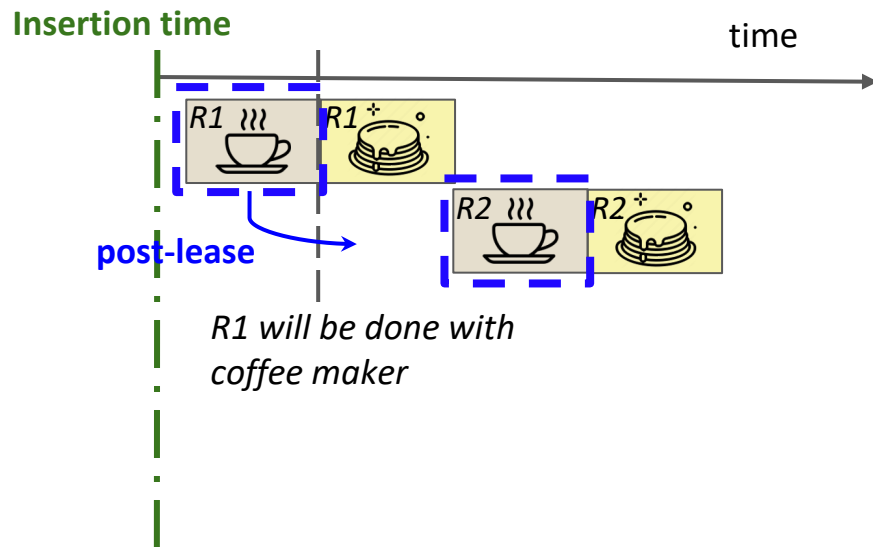
- Routines can concurrently execute *without violating some serial order*.
- Each routine holds the **locks** for devices it touches (but can **lease the lock**).



# Eventual Visibility (EV) - Post-Lease

Post-lease:

- If a routine is done with a device  $D$ , it can post-lease  $D$ 's lock to another routine.



Serial order:  
lessor  $\rightarrow$  lessee  
(  $R1 \rightarrow R2$  )

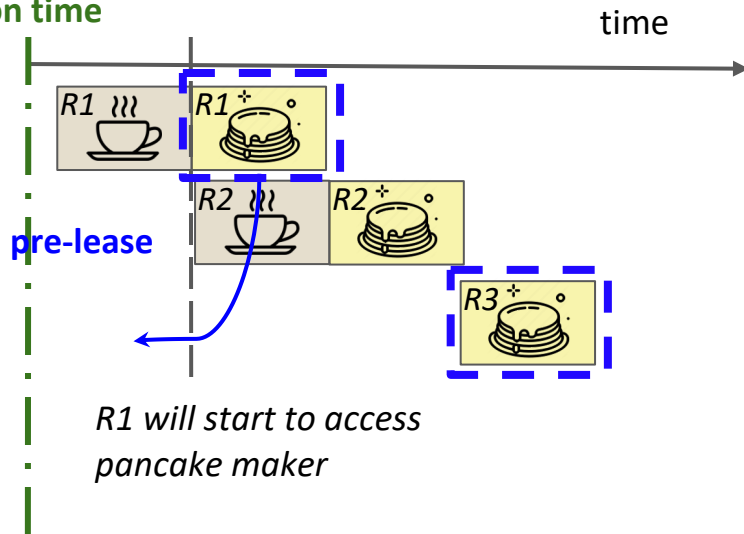


# Eventual Visibility (EV) - Pre-Lease

Pre-lease:

- If a routine has acquired the lock but not accessed a device  $D$ , it can pre-lease  $D$ 's lock to another routine.

Insertion time



Serial order:  
lessee  $\rightarrow$  lessor  
(  $R3 \rightarrow R1$  )



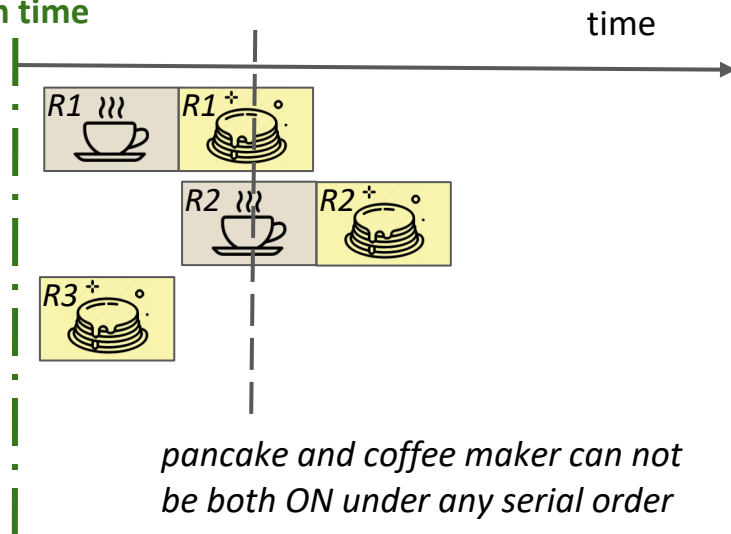
# Eventual Visibility (EV)

EV finishes routine

- with *short wait* and provides *serial equivalence*
- with higher *temporary incongruence*: intermediate state is not serially equivalent

Finish in 3 time units

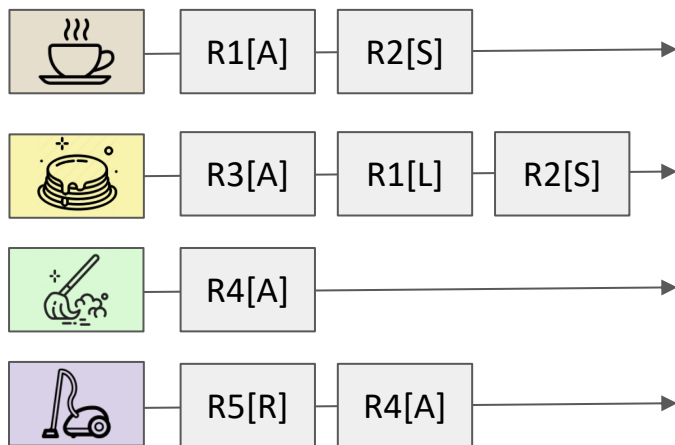
Insertion time





# Eventual Visibility (EV) - Lineage Table

*Lineage Table:* SafeHome's plan of which routine will access which device.



[A]: Get lock Access  
[S]: Routine Scheduled  
[L]: Lock Leased out  
[R]: Lock Released

## *Scheduling plan placement:*

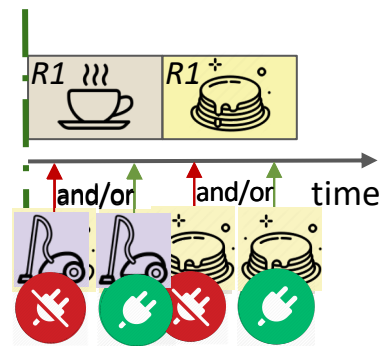
- Placed when routine is triggered
- Use *backtracking* for valid placement
- Explore two other policies (FCFS, JiT)

# Failure Serialization and Rollback

Device might *fail*:

- *Rollback*? Try to *serialize* the failure/restart event!
- If the failed device is not touched by the routine:
  - *Arbitrary* Serial Equivalence order
- If device fails/restarts after the last touch:
  - *Routine* → *Fail/Restart* Serial Equivalence order
- If device fails/restarts before the first touch:
  - *Fail/Restart* → *Routine* Serial Equivalence order
- If device fails/restarts during the touch:
  - *Rollback* routine

Start Execution



*Failure* → *Restart* → *R1*



# SafeHome Implementation

## Implementation

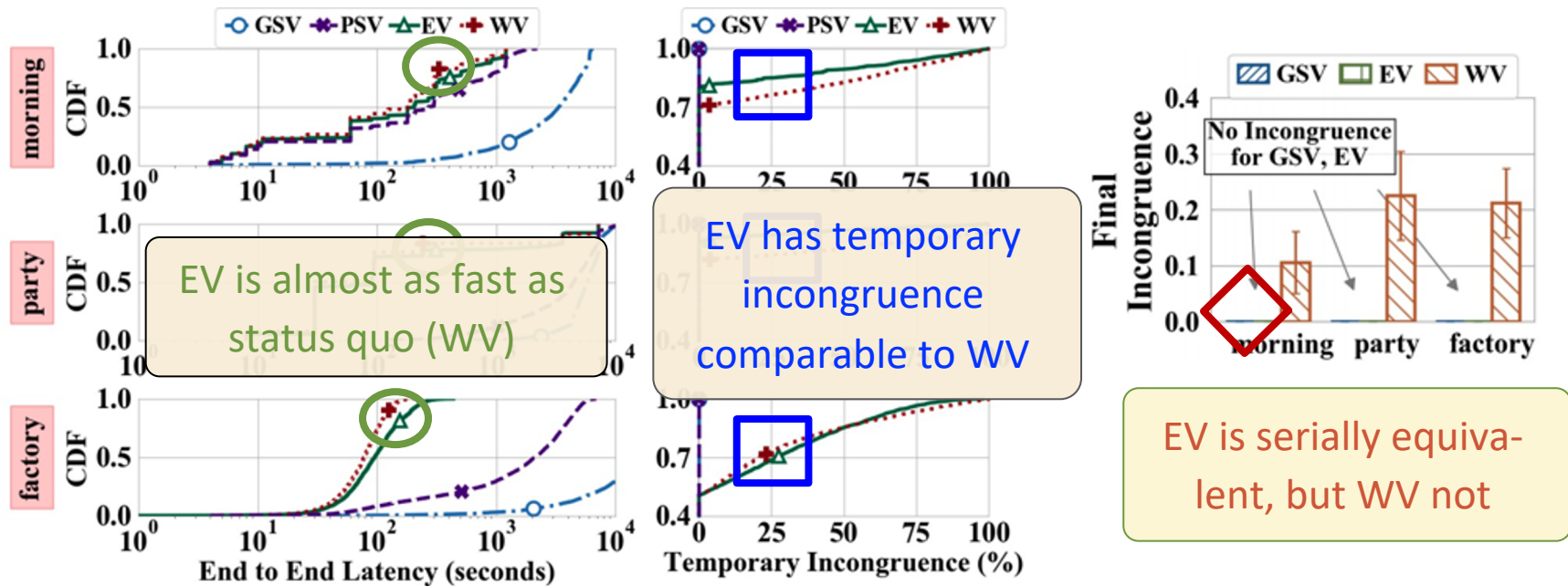
- ~2k line of *Java* code
- Support *long running routine* expression (JSON)
- Popular Smart Device *integration* (TP-link, Wemo)

## Experiment Setup

- Deployment & Simulation
- Real-world Benchmark
  - Derived from *IoT Bench Test Suite*
  - *Morning, Party, Factory* Scenario
- Workload-Driven
  - Average of *500k* runs



# Real-World Benchmark



*Temporary Incongruence:* the ratio of time when *intermediate* state is not serially equivalent.

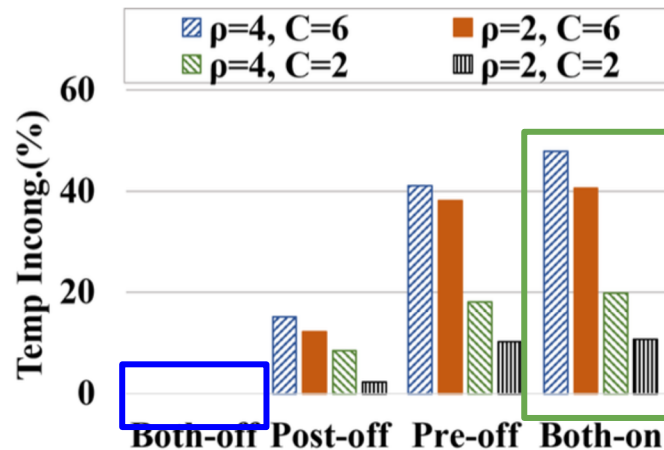
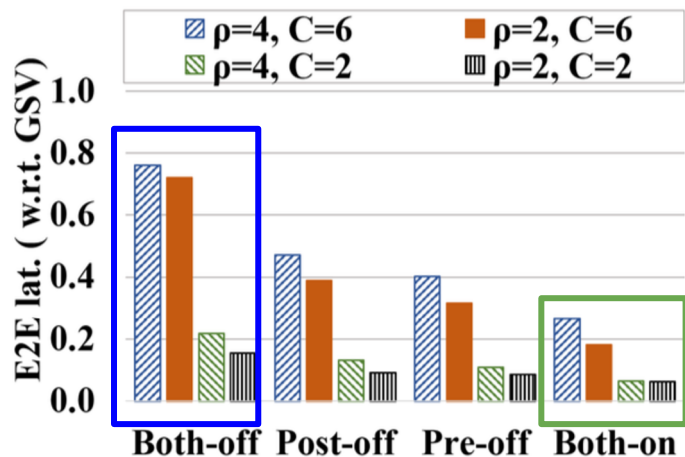
*Final Incongruence:* the ratio of runs that *end up* in an incongruent state.



# Workload Evaluation -- Pre/Post-Lease

High Latency, Zero Temporary Incongruence

Low Latency, High Temporary Incongruence



*Pre/Post leases reduce the E2E latency (user-facing metrics) with the cost of Temporary Incongruence*



# Takeaways

- Safehome is a first step to provide *reliability* from routine execution level
- SafeHome provides four *Visibility Models* (WV, EV, PSV, and GSV)
- *Eventual Visibility* (EV) model provides the best of both worlds, with:
  - Good user-facing *responsiveness* (0 - 23.1%)
  - Strongest *end state congruence* equivalent guarantee (as GSV)
- Lock-leasing *improves latency* by 1.5X - 4X

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